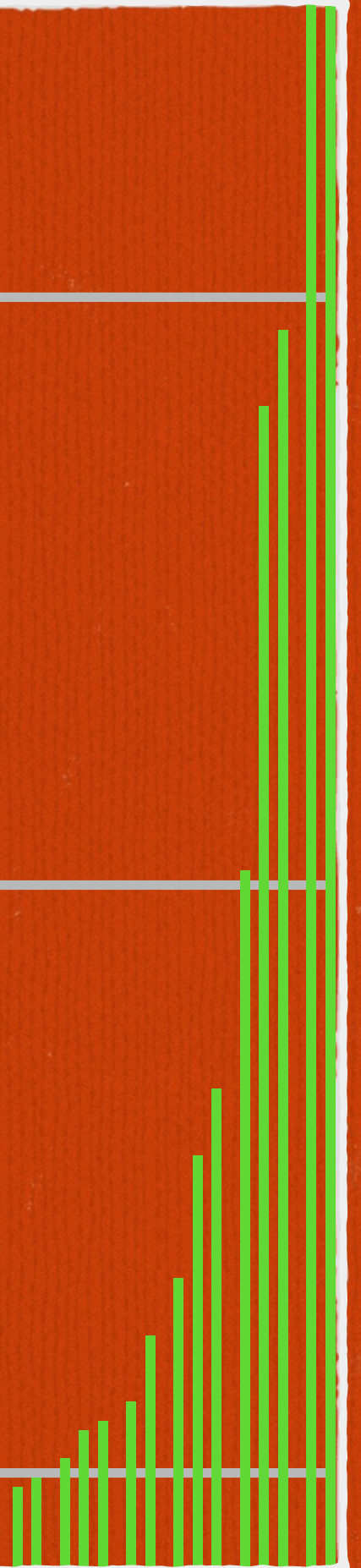


15.000.000

10.000.000

5.000.000



30 Years of CPS

The Rexx Clauses-per-second benchmark

The Standard by MFC

```
CPS EXEC A1          At: " /* Aver".. 2 changes

/*-----*/
rexxcps=2.1 /* REXXCPS version; quotable only if code unchanged */

/* Adjust these counts so run-time is about one second */
count=100 /* Repetition count */
averaging=100 /* Averaging-over count */

tracevar='Off' /* Trace setting (for development use) */

signal on novalue
parse source source 1 system .
parse version version

say '----- REXXCPS' rexxcps '-- Measuring REXX clauses/second -----'
say ' REXX version is:' version
say ' System is:' system
say ' Averaging:' averaging 'measures of' count 'iterations'

/* Calibrate for the empty do-loop */
empty=0
do i=1 to averaging
  call time 'R'
  do count; end
  empty=time('R')+empty
end
empty=empty/averaging

noterm=(system='CMS'); if pos('0',tracevar)=1 then noterm=0
if noterm then do
  say 'Calibration (empty DO):' empty 'secs (average of' averaging')'
  say 'Spooling trace NOTERM'
  'CP SPOOL CON * START NOTERM'; 'CP CLOSE CON PUR'
end

PF2 Openc1 PF3 Quit PF4 Copy PF5 Move PF6 ?
PF8 Down PF9 = PF10 Select PF11 Splitj PF12 Focus
```

- Measures Clauses Per Second
 - A clause is ~ a line of instructions
- History of measurements since 1989
- Product of analysis of 1000's of lines of real applications
- Multiplatform Classic Rexx

What does it do?

- Two loops**
 - One for calibration of an (almost) empty loop**
 - One that does the work**
- Can specify # of averaging and measuring loops**
- Total execution should be 1 sec elapsed time for dependable results**

Calibrating the empty loop

```
CPS      EXEC      B1  V 130  Trunc=130 Size=141 Line=0 Col=1 Alt=0
====>
|...+...1...+...2...+...3...+...4...+...5...+...6...+...
00000 * * * Top of File * * *
00001 ----- 39 line(s) not displayed -----
00040 /* Calibrate for the empty do-loop */
00041 empty=0
00042 do i=1 to averaging
00043   call time 'R'
00044   do count; end
00045   empty=time('R')+empty
00046   end
00047 empty=empty/averaging
00048
00049 noterm=(system='CMS'); if pos('0',tracevar)=1 then noterm=0
00050 if noterm then do
00051   say 'Calibration (empty DO):' empty 'secs (average of' averaging')'
00052   say 'Spooling trace NOTERM'
00053   'CP SPOOL CON * START NOTERM'; 'CP CLOSE CON PUR'
00054   end
00055 ----- 87 line(s) not displayed -----
00142 * * * End of File * * *
```

```

CPS      EXEC      B1  V 130  Trunc=130 Size=141 Line=0 Col=1
====>
|...+....1....+....2....+....3....+....4....+....5....+
00000 * * * Top of File * * *
00001 ----- 55 line(s) not displayed -----
00056 /* Now the true timer loop .. average timing again */
00057 full=0
00058 do i=1 to averaging
00059   trace value tracevar
00060   call time 'R'
00061   do count;
00062     /* ----- This is first of the 1000 clauses ----- */
00063     flag=0; p0='b'
00064     do loop=1 to 14
00065       /* This is the "block" comment in loop */
00066       key1='Key Bee'
00067       acompound.key1.loop=substr(1234"5678",6,2)
00068       if flag=acompound.key1.loop then say 'Failed1'
00069       do j=1.1 to 2.2 by 1.1 /* executed 28 times */
00070         if j>acompound.key1.loop then say 'Failed2'
00071         if 17<length(j)-1 then say 'Failed3'
00072         if j='foobar' then say 'Failed4'
00073         if substr(1234,1,1)=9 then say 'Failed5'
00074         if word(key1,1)='?' then say 'Failed6'
00075         if j<5 then do /* This path taken */
00076           acompound.key1.loop=acompound.key1.loop+1
00077           if j=2 then leave
00078         end
00079         iterate
00080       end /* j */
00081     avar.=1.0''loop
00082     select
00083       when flag='string' then say 'FailedS1'
00084       when avar.flag.2=0 then say 'FailedS2'
00085       when flag=5+99.7 then say 'FailedS3'
00086       when flag then avar.1.2=avar.1.2*1.1
00087       when flag==0 then flag=0
00088     end

```

The timer loop

```

CPS      EXEC      B1  V 130  Trunc=130 Size=141 Line=88 Col=1 Alt=0
====>
|...+....1....+....2....+....3....+....4....+....5....+....6....+....7.
00088   end
00089   if 1 then flag=1
00090   select
00091     when flag=='ring' then say 'FailedT1'
00092     when avar.flag.3=0 then say 'FailedT2'
00093     when flag then avar.1.2=avar.1.2*1.1
00094     when flag==0 then flag=1
00095   end
00096   parse value 'Foo Bar' with v1 +5 v2 .
00097   trace value trace(); address value address()
00098   call subroutine 'with' 2 'args', '(This is the second)'1'1
00099   rc='This is an awfully boring program'; parse var rc p1 (p0) p5
00100   rc='is an awfully boring program This'; parse var rc p2 (p0) p6
00101   rc='an awfully boring program This is'; parse var rc p3 (p0) p7
00102   rc='awfully boring program This is an'; parse var rc p4 (p0) p8
00103   end loop
00104   /* ----- This is last of the 1000 clauses ----- */
00105   end
00106   full=time('R')+full
00107   trace off
00108   end
00109   full=full/averaging
00110 ----- 32 line(s) not displayed -----
00142 * * * End of File * * *

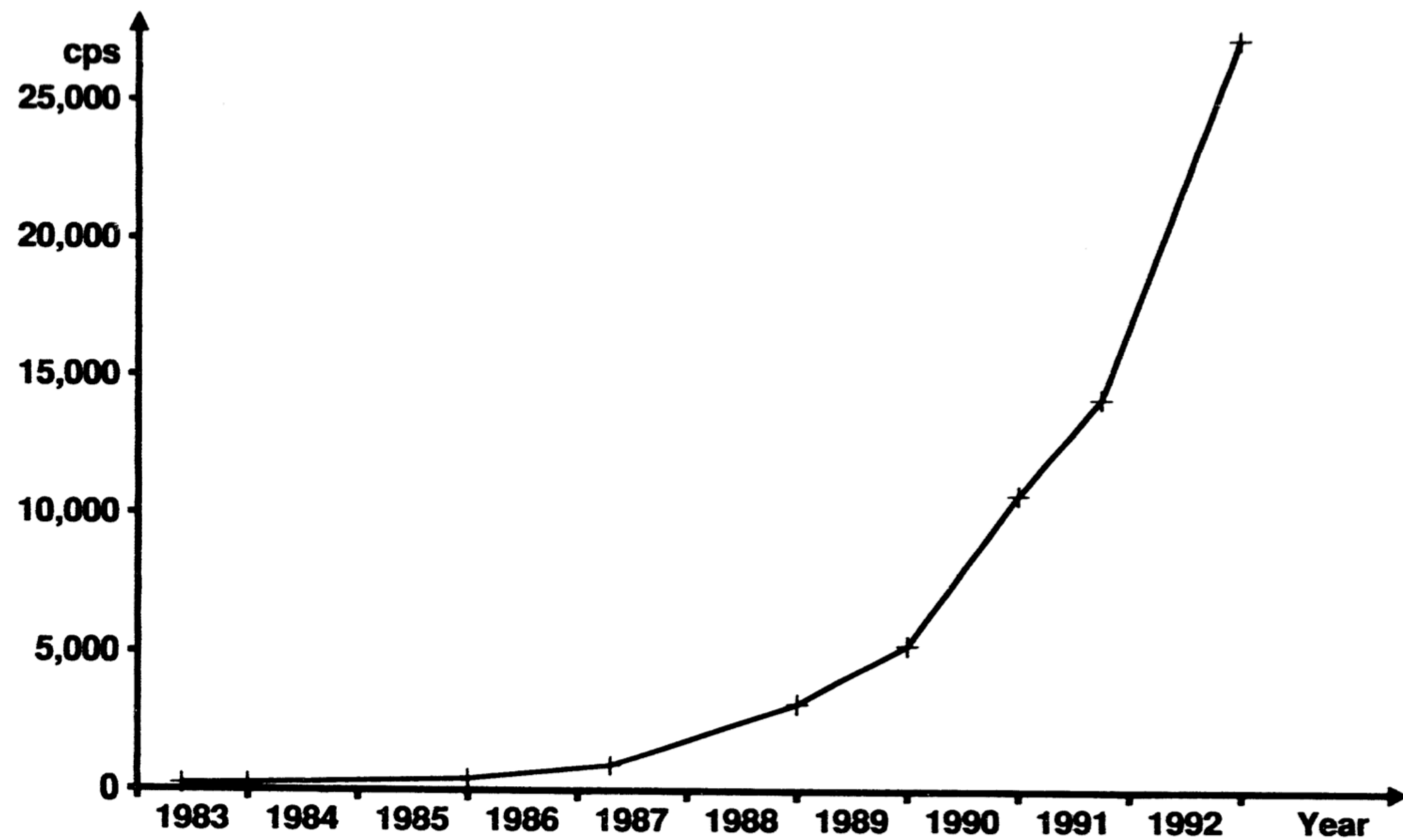
```

Hardware speed and REXX

As hardware speed increases, REXX is being used for a wider set of applications. Some informal figures:

- ◆ x86 systems—now over 27,000 REXX clauses per second (486/66)
- ◆ RISC systems—over 42,000 REXX cps (same interpreter)
- ◆ Mainframe systems—over 90,000 REXX cps
- ◆ REXX Compiler/370—up to 465,000 REXX cps

REXX Clauses per Second—x86 platform



13
13 May 1993

- 4 -

Mike Cowlshaw

WE APPEAR TO HAVE NUMBERS STARTING 1983

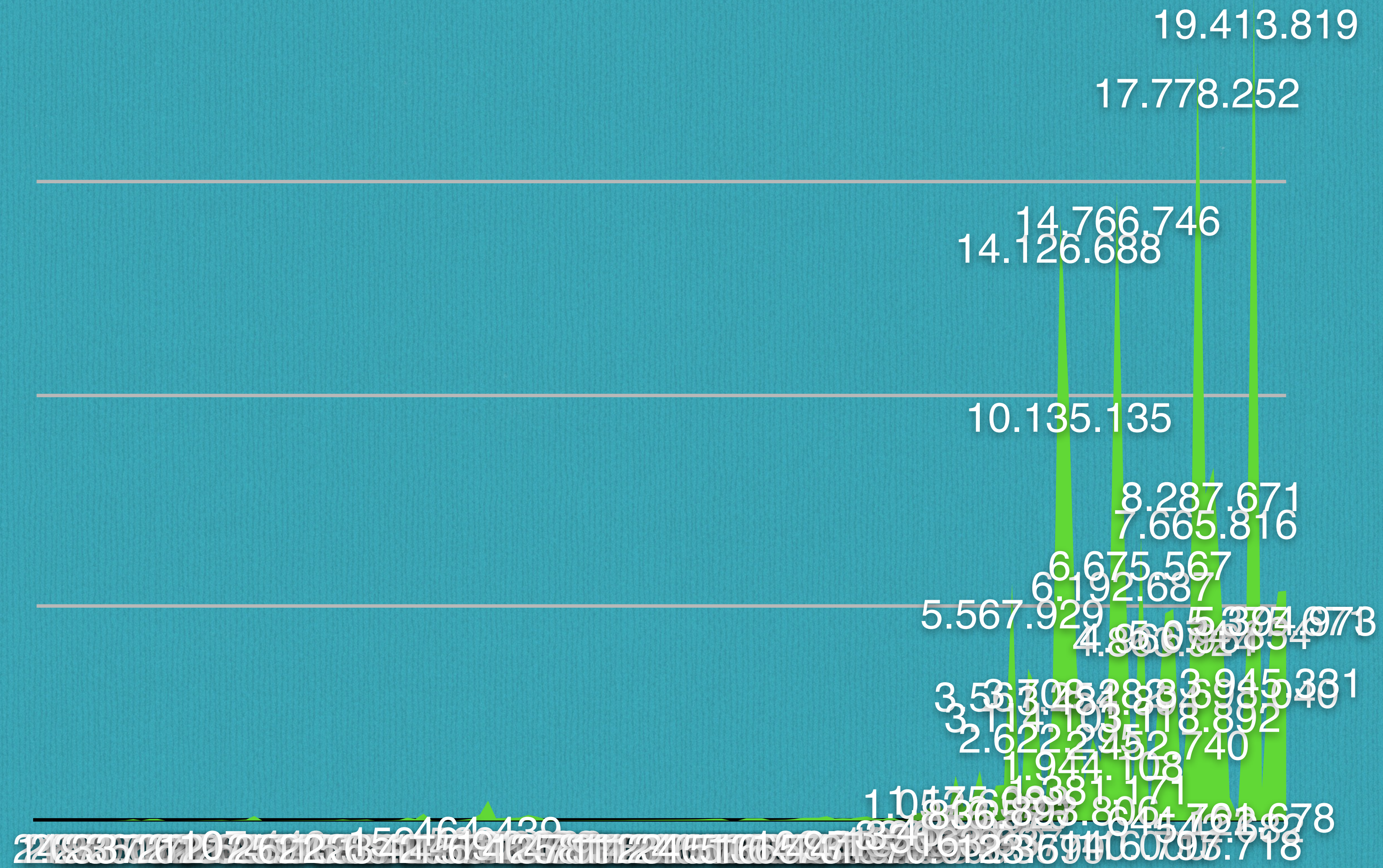
20.000.000

15.000.000

10.000.000

5.000.000

0



1992

2015

- But which is the 'standard' standard?
- Every interpreter delivers one
- All ever so slightly different
 - Brexx and ooRexx cannot run each other's rexxcps
 - When fixed, they run their own cps faster than the other's
- Canonical version on speleotrove.com

Influence of compiler options

Interpreter	Version	Architecture	OS	Virtualization	Compiler	Compiler Version	Compiler options	CPS score	Hardware	CPU	Date	Measures/ Iterations
brexx	2.1.9	s390x	Red Hat Enterprise Linux Server 7.5 (Maipo)	z/VM	gcc	4.8.5	"-O3 fno-stack-protector	6.961.707	IBM LinuxOne Rockhopper (Z13s)		2019-07-04	100/100
brexx	2.1.9	s390x	Red Hat Enterprise Linux Server 7.5 (Maipo)	z/VM	gcc	4.8.5	+ -mzarch	8.502.614	IBM LinuxOne Rockhopper (Z13s)		2019-07-04	100/100
brexx	2.1.9	s390x	Red Hat Enterprise Linux Server 7.5 (Maipo)	z/VM	gcc	4.8.5	+ "-m-hard-float" "-m-hard-ldfp"	8.516.123	IBM LinuxOne Rockhopper (Z13s)		2019-07-04	100/100
brexx	2.1.9	s390x	Red Hat Enterprise Linux Server 7.5 (Maipo)	z/VM	gcc	4.8.5	+ "-march=z13"	8.567.324	IBM LinuxOne Rockhopper (Z13s)		2019-07-04	100/100
brexx	2.1.9	s390x	Red Hat Enterprise Linux Server 7.5 (Maipo)	z/VM	gcc	4.8.5	static executable	8.827.065	IBM LinuxOne Rockhopper (Z13s)		2019-07-04	100/100
REXX370 4.02 01 Dec 1998	4.02	z/Arch	CMS 6.4.0	z/VM				3.246.319	IBM LinuxOne Rockhopper (Z13s)		2018-11-15	100/100
brexx	2.1.9	x86_64	Ubuntu 18.04.2 LTS	no	gcc	7.4.0		19.620.712	Dell Latitude 5480	Intel Corei7-600U	2019-07-04	100/100

ARCHLVL

FP/DFP

Current numbers on x86_64 (Intel I9)

BREXX	21104425
ooRexx	13533448
Regina	8238630
Jaxx	9278788

```

Architecture:      x86_64
CPU op-mode(s):    32-bit, 64-bit
Byte Order:        Little Endian
CPU(s):            4
On-line CPU(s) list: 0-3
Thread(s) per core: 2
Core(s) per socket: 2
Socket(s):         1
NUMA node(s):     1
Vendor ID:         GenuineIntel
CPU family:        6
Model:             142
Model name:        Intel(R) Core(TM) i7-7567U CPU @ 3.50GHz
Stepping:          9
CPU MHz:           1000.053
CPU max MHz:       4000.0000
CPU min MHz:       400.0000
BogoMIPS:          7008.00
Virtualization:    VT-x
L1d cache:         32K
L1i cache:         32K
L2 cache:          256K
L3 cache:          4096K
NUMA node0 CPU(s): 0-3
Flags:              fpu vme de pse tsc msr pae mce cx8 apic sep mtrr pge mca c
constant_tsc art arch_perfmon pebs bts rep_good nopl xtopology nonstop_tsc
fma cx16 xtpr pdc_m pcid sse4_1 sse4_2 x2apic movbe popcnt tsc_deadline_
ibrs ibpb stibp tpr_shadow vnmi flexpriority ept vpid fsgsbase tsc_adjust bmi
xsaves dtherm ida arat pln pts hwp hwp_notify hwp_act_window hwp_epp m

```

Current numbers on ARM 64 (aarch64)

BREXX	2648138
ooRexx	1831838
Regina	2062977

```
rvjansen@jetson:~/data/brex/src$ lscpu
```

```
Architecture: aarch64
```

```
Byte Order: Little Endian
```

```
CPU(s): 4
```

```
On-line CPU(s) list: 0-3
```

```
Thread(s) per core: 1
```

```
Core(s) per socket: 4
```

```
Socket(s): 1
```

```
Vendor ID: ARM
```

```
Model: 1
```

```
Model name: Cortex-A57
```

```
Stepping: r1p1
```

```
CPU max MHz: 1428.0000
```

```
CPU min MHz: 102.0000
```

```
BogoMIPS: 38.40
```

```
L1d cache: 32K
```

```
L1i cache: 48K
```

```
L2 cache: 2048K
```

```
Flags: fp asimd evtstrm aes pmull sha1 sha2 crc32
```

Current numbers on ARM 32 (armv7l)

BREXX	2816483
ooRexx	1914118
Regina	2043062

pi@kleene:~/apps \$ lscpu

Architecture: armv7l

Byte Order: Little Endian

CPU(s): 4

On-line CPU(s) list: 0-3

Thread(s) per core: 1

Core(s) per socket: 4

Socket(s): 1

Vendor ID: ARM

Model: 3

Model name: Cortex-A72

Stepping: r0p3

CPU max MHz: 1500.0000

CPU min MHz: 600.0000

BogoMIPS: 108.00

Flags: half thumb fastmult vfp edsp neon vfpv3 tls vfpv4 idiva ic

Why is ooRexx slower?

- Interpreter api
- Stems are slower than array objects
-

Why is BREXX faster?

- Limited numeric precision
- Not an exact implementation of the standard

CRX

- Written in Intel assembler (masm) for DOS and Windows, very fast**
- Its purpose was to verify the ANSI (ISO) standard**
- Parts of it generated from the grammar in that document**
- Made with knowledge of Intel processor cache and pipelines**

Parallel Rexx

- Maybe the future
- Nvidia and Intel have competing 'AI' processors
- CUDA in video cards exists for several years now, CUDA C++ can - with your assistance, that is - parallelize selected loops to execute hundreds times faster, consistent with the number of GPU's available
- To be faster we cannot really trust the CPUs itself
- Experiments would be very interesting

NetRexx CPS

There never was one

Or was there?

NetRexx



Links

- [Learn more ...](#)
- [NetRexx Tools](#)
- [NetRexx Forum](#)
- [Rexx Language Association](#)

Welcome to the NetRexx Programming Language!

NetRexx is a general-purpose programming language inspired by two very different programming languages: Rexx and Java. In this respect it follows Rexx closely, with many of the concepts and most of the syntax of Rexx. From Java it derives static typing, binary arithmetic, the object model, and exception handling. Originally a product from the IBM Hursley Software Lab, NetRexx has always been free software. It is an alternative language for the Java Virtual Machine (JVM).

Why NetRexx?

- NetRexx makes programming easy, and fun again
- The Rexx type, combining numeric and string processing in one class
- Unlimited precision arithmetic built into the language
- Interpret your code or compile to JVM .class files
- Lightning fast performance compared to other JVM scripting languages
- Seamless integration to all JVM libraries
- Built-in luxurious parsing and tracing

Hursley Labs

Hursley, located near Winchester in the UK, is the place where many famous products were developed. It was here that the [qtime](#) program, one of the first-ever Rexx programs, dating from 1979. This is the first Rexx program that NetRexx listens to both *center* and *centre* method spellings.

MFC on NetRexx

The seamless integration of types into what was previously an essentially typeless language is a strong typing while preserving the ease of use and speed of development that Rexx provides. The addition of these capabilities to the Rexx language is a single language that has both the Rexx strengths of robustness, good efficiency, portability, and security for application development.
-- Mike Cowlishaw

rexxcps.nrx

- From 1996, by Mike Cowlshaw, but never released until now**
- Optimizing compilers optimize a lot of code away, must make sense of the numbers**
- In NetRexx 3.08, there is an experimental version released**
 - Uses Nanotime instead of millisecs**
 - Currently evaluated on different architectures to see what actually is executed**

- NetRexx programs have a complex runtime**
- NetRexx is translated to Java source, which is compiled to bytecode, which is interpreted (sometimes) but mostly compiled just-in-time to native code**
- Optimization takes place on different levels, each can throw out code**
- Reaches extremely high numbers**

Current rexxcps.nrx

Aarch64	2,407E+15	2_407_464_380_676
IA 86_64 I9	1,1489E+16	11_488_524_444_376
IA 86_64 I7	9,861E+15	9_861_203_229_267

How do we see what is executed?

List the generated java

List the generated bytecode

List the generated assembler

```
File Edit Options Buffers Tools Index NetRexx Help
say '          Averaging:' averaging 'measures of' count 'iteration'

/* ----- Calibrate for the empty loop ----- */
empty=long 0
loop for averaging
  start=System.nanoTime()
  loop for count; end
  empty=empty+System.nanoTime()-start
end
fempty=empty/averaging

/* ----- Now the true timer loop .. average timing again ----- */
full=long 0
loop for averaging
  start=System.nanoTime()
  loop for count
    /* ----- This is first of the 1000 clauses ----- */
    flag=0; p0='b'
    loop lvar=1 to 14
      /* This is the "block" comment in lvar */
      key1='Key Bee'; acompound=''
      acompound[key1,lvar]=(1234"5678").substr(6,2)
      if flag=acompound[key1,lvar] then say 'Failed1'
      loop j=1.1 to 2.2 by 1.1 /* executed 28 times */
        if j>acompound[key1,lvar] then say 'Failed2'
        if 17<key1.length-1 then say 'Failed3'
        if j='foobar' then say 'Failed4'
        if key1.substr(1,1)=9 then say 'Failed5'
        if key1.word(1)='?' then say 'Failed6'
        if j<5 then do /* This path taken */
          acompound[key1,lvar]=acompound[key1,lvar]+1
          if j=2 then leave /* never */
        end
      end
    iterate
  end /* j */
  avar=1.0''lvar
select
```

.nrx yields .java

We can verify a 1:1 relationship between the nrx and the java sourcecode

```
ssh
File Edit Options Buffers Tools Java Help
{int $5=count.OpPlus($1).toint();for(; $5>0;$5--){
}
}
empty=new netrexx.lang.Rexx(empty).OpAdd($1,new netrexx.lang.Rexx(java.lang.
(start)).tolong());
}

empty=new netrexx.lang.Rexx(empty).OpDiv($1,averaging);
null=new netrexx.lang.Rexx(0).tolong();
int $6=averaging.OpPlus($1).toint();for(; $6>0;$6--){
start=java.lang.System.nanoTime();
{int $7=count.OpPlus($1).toint();for(; $7>0;$7--){
flag=new netrexx.lang.Rexx((byte)0);
p0=new netrexx.lang.Rexx('b');
{lvar=new netrexx.lang.Rexx((byte)1);lvar:for(;lvar.OpLtEq($1,$07);lvar=lva
key1=netrexx.lang.Rexx.toRexx("Key Bee");
acomound=netrexx.lang.Rexx.toRexx("");
acomound.getnode(key1).leaf.getnode(lvar).leaf=($08.OpCc($1,$09)).substr(
xx((byte)2));
if (flag.OpEq($1,acomound.getnode(key1).leaf.getnode(lvar).leaf))
netrexx.lang.RexxIO.Say("Failed1");
{netrexx.lang.Rexx $8=$010;j=$010.OpPlus($1);j:for(;j.OpLtEq($1,$011);j=j.
if (j.OpGt($1,acomound.getnode(key1).leaf.getnode(lvar).leaf))
netrexx.lang.RexxIO.Say("Failed2");
if ($013.OpLt($1,(key1.length()).OpSub($1,$012)))
netrexx.lang.RexxIO.Say("Failed3");
if (j.OpEq($1,$014))
netrexx.lang.RexxIO.Say("Failed4");
if ((key1.substr(new netrexx.lang.Rexx((byte)1),new netrexx.lang.Rexx((by
netrexx.lang.RexxIO.Say("Failed5");
if ((key1.word(new netrexx.lang.Rexx((byte)1))).OpEq($1,$016))
netrexx.lang.RexxIO.Say("Failed6");
if (j.OpLt($1,$017))
{
acomound.getnode(key1).leaf.getnode(lvar).leaf=(acomound.getnode(key1
if (j.OpEq($1,$018))
break j;
}:----F1 rexxcps.java 40% L107 (Java/l Abbrev) -----
```

.java yields .class

We can verify that all bytecode for the benchmark has been generated

```
540: new          #1          // class netrexx/lang/F
543: dup
544: iconst_2
545: invokespecial #7          // Method netrexx/lang/
548: invokevirtual #39         // Method netrexx/lang/
</lang/Rexx;
551: putfield     #35         // Field netrexx/lang/F
554: aload       15
556: getstatic   #11         // Field $1:Lnetrexx/la
559: aload       19
561: aload       18
563: invokevirtual #34         // Method netrexx/lang/
566: getfield    #35         // Field netrexx/lang/F
569: aload       17
571: invokevirtual #34         // Method netrexx/lang/
574: getfield    #35         // Field netrexx/lang/F
577: invokevirtual #40         // Method netrexx/lang/
580: ifeq        589
583: ldc         #41         // String Failed1
585: invokestatic #42         // Method netrexx/lang/
588: pop
589: getstatic   #43         // Field $010:Lnetrexx/
592: astore     39
594: getstatic   #43         // Field $010:Lnetrexx/
597: getstatic   #11         // Field $1:Lnetrexx/la
600: invokevirtual #22         // Method netrexx/lang/
603: astore     20
605: aload       20
607: getstatic   #11         // Field $1:Lnetrexx/la
610: getstatic   #44         // Field $011:Lnetrexx/
613: invokevirtual #31         // Method netrexx/lang/
616: ifeq        867
619: aload       20
621: getstatic   #11         // Field $1:Lnetrexx/la
624: aload       19
626: aload       18
628: invokevirtual #34         // Method netrexx/lang/
-UU-:----F1 rexcps.b 19% L344 (Fundamental) -----
```


The Empty Averaging Loop in Java

```
ssh
File Edit Options Buffers Tools Java Help
count=new netrexx.lang.Rexx((byte)100);
averaging=new netrexx.lang.Rexx((byte)100);
rexxcps=netrexx.lang.Rexx.toRexx("2.1n");
netrexx.lang.RexxIO.Say((netrexx.lang.Rexx.toRexx("----- REXXCPS").OpCcbblank($1, rexxcps)).OpCcbblank($1, netrexx.lang.Rexx.toR\
exx("-- Measuring NetRexx clauses/second -----")));
netrexx.lang.RexxIO.Say(netrexx.lang.Rexx.toRexx(" NetRexx version is:").OpCcbblank($1, version));
netrexx.lang.RexxIO.Say(netrexx.lang.Rexx.toRexx(" System is:").OpCcbblank($1, opsystem));
netrexx.lang.RexxIO.Say(((netrexx.lang.Rexx.toRexx(" Averaging:").OpCcbblank($1, averaging)).OpCcbblank($1, netrexx.la\
ng.Rexx.toRexx("measures of"))).OpCcbblank($1, count)).OpCcbblank($1, $06));
empty=new netrexx.lang.Rexx(0).tolong();
{int $4=averaging.OpPlus($1).toint();for(; $4>0; $4--){
  start=java.lang.System.nanoTime();
  {int $5=count.OpPlus($1).toint();for(; $5>0; $5--){
    }
  }
  empty=new netrexx.lang.Rexx(empty).OpAdd($1, new netrexx.lang.Rexx(java.lang.System.nanoTime())).OpSub($1, new netrexx.lang.R\
exx(start)).tolong();
}
fempty=new netrexx.lang.Rexx(empty).OpDiv($1, averaging);
full=new netrexx.lang.Rexx(0).tolong();
{int $6=averaging.OpPlus($1).toint();for(; $6>0; $6--){
  start=java.lang.System.nanoTime();
  {int $7=count.OpPlus($1).toint();for(; $7>0; $7--){
    flag=new netrexx.lang.Rexx((byte)0);
    p0=new netrexx.lang.Rexx('b');
    {lvar=new netrexx.lang.Rexx((byte)1);lvar:for(; lvar.OpLtEq($1, $07); lvar=lvar.OpAdd($1, new netrexx.lang.Rexx(1))){
      key1=netrexx.lang.Rexx.toRexx("Key Bee");
      acompound=netrexx.lang.Rexx.toRexx("");
      acompound.getnode(key1).leaf.getnode(lvar).leaf=( $08.OpCc($1, $09)).substr(new netrexx.lang.Rexx((byte)6), new netrexx.lang\
.Rexx((byte)2));
      if (flag.OpEq($1, acompound.getnode(key1).leaf.getnode(lvar).leaf))
        netrexx.lang.RexxIO.Say("Failed1");
      {netrexx.lang.Rexx $8=$010; j=$010.OpPlus($1); j:for(; j.OpLtEq($1, $011); j=j.OpAdd($1, $8)){
        if (j.OpGt($1, acompound.getnode(key1).leaf.getnode(lvar).leaf))
          netrexx.lang.RexxIO.Say("Failed2");
        if ($013.OpLt($1, (key1.length()).OpSub($1, $012))
          netrexx.lang.RexxIO.Say("Failed3");
      }
    }
  }
}
-UU-:----F1 rexxcps.java 33% L108 (Java/l Abbrev) -----
```

Generated Assembler

- With a special dll we can see what the HotSpot compiler generates as assembly
- We can also see and influence which parts will be compiled and which parts will be interpreted

Parallelism

- Rexxcps.nrx also is a single tasking program**
- As java can also be parallelized, the future might hold change**
- Pipelines in NetRexx, for example are already multitasked over all available processors**

```
→ test git:(master) x time pipe 'literal aap noot mies | split | reverse | a: locate seim ? a: | console'  
paa  
toon  
java org.netrexx.njpipes.pipes.runner 2.06s user 0.12s system 249% cpu 0.870 total  
→ test git:(master) x █
```

Example of Parallelism

2.067 seconds are spent with 249% cpu, which makes for an elapsed time of 0.87 seconds